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HAL Id: halshs-00186905
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Submitted on 13 Nov 2007

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Dominance of capacities by $k$-additive belief functions

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Abstract

In this paper we deal with the set of $k$-additive belief functions dominating a given capacity. We follow the line introduced by Chateauneuf and Jaffray for dominating probabilities and continued by Grabisch for general $k$-additive measures. First, we show that the conditions for the general $k$-additive case lead to a very wide class of functions and this makes that the properties obtained for probabilities are no longer valid. On the other hand, we show that these conditions cannot be improved. We solve this situation by imposing additional constraints on the dominating functions. Then, we consider the more restrictive case of $k$-additive belief functions. In this case, a similar result with stronger conditions is proved. Although better, this result is not completely satisfactory and, as before, the conditions cannot be strengthened. However, when the initial capacity is a belief function, we find a subfamily of the set of dominating $k$-additive belief functions from which it is possible to derive any other dominant $k$-additive belief function, and such that the conditions are even more restrictive, obtaining the natural extension of the result for probabilities. Finally, we apply these results in the fields of Social Welfare Theory and Decision Under Risk.

Keywords: Linear programming, decision analysis, capacity, dominance, $k$-additivity, belief functions.

1 Introduction

The problem of finding the set of probabilities dominating a capacity has been extensively studied; for instance, Dempster [3] and Shafer [25] have proposed a representation of uncertainty based on a “lower probability” or “degree of belief”, respectively, to every event. Their model needs a lower probability function, usually non-additive but having a weaker property: monotonicity of order $k$, for all $k$, i.e. it is a belief function [25]. This requirement is perfectly justified in some situations (see [3]). The general form of lower probabilities has been studied by several authors (see e.g. [31, 32]).

Moreover, in many decision problems, in which we have not enough information, decision makers often feel that they are only able to assign an interval value for the probability of events. In other words, the real probability distribution is unknown, but there exists a set of probabilities compatible with the available information. Let us call this set of all compatible additive probabilities $P_1$ and let us define $\mu := \inf_{P \in P_1} P$; then, $\mu$ is a capacity (but not necessarily a belief function [30]); $\mu$ is called “coherent lower probability”, and it is the natural “lower

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*This paper is a revised and extended version of [19], presented at IFSA2003 conference.
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probability function”. Of course, if \( P \) is a probability distribution dominating \( \mu \), it is clear that \( E_{P'}(f) \geq C_{\mu}(f) \), for any function \( f \), where \( C_{\mu} \) represents Choquet integral \([2]\) (function \( f \) plays the role of an action in the Decision problem \([27]\)). Chateauneuf and Jaffray use in \([1]\) this fact and that \( \mu \leq P, \forall P \in \mathcal{P} \), to obtain an easy method to compute \( \inf_{P \in \mathcal{P}} E_{P}(f) \). This method is based on the set of all probability distributions dominating \( \mu \). The same can be done for obtaining an upper bound.

In our case, we try to extend these results for \( k \)-additive measures. There are some cases where this could be useful. First, suppose a situation that can be modelled by a \( k \)-additive measure (an axiomatic approach to this situation can be found in \([17, 22]\)), but where our information does not allow us to completely determine the measure. Then, we have to work with a set of compatible \( k \)-additive measures (let us call it \( U_{k} \)). A second example is the identification of a capacity in a practical situation. It can be proved that sometimes it is not possible to determine a single solution but a set of \( k \)-additive measures, all equally suitable \([16]\); it can also be proved that the set of all suitable measures is a convex set and consequently, the measure for an event \( A \subseteq X \) can take an interval of possible values (a deeper study about the unicity of solution and the structure of the set of solutions can be found in \([20]\)). If \( \mu = \inf_{\mu_{k} \in U_{k}} \mu_{k} \), then \( \mu \) is clearly a capacity \([27]\), and we have the inequality \( C_{\mu_{k}}(f) \geq C_{\mu}(f) \), for any \( k \)-additive measure \( \mu'_{k} \) dominating \( \mu \) and for any function \( f \). Therefore, it seems interesting to find the set of all \( k \)-additive fuzzy measures dominating \( \mu \) in order to try to translate the results in \([1]\). However, we will see later that the results for the general \( k \)-additive case are too weak to deal with. Thus, we will focus on the special case of belief functions.

The paper is organized as follows: In next section we introduce the basic concepts on fuzzy measures that we will use in the paper. In Section 3 we recall the results on dominating probabilities and \( k \)-additive measures. We also show that the conditions for the \( k \)-additive case are too general in order to maintain the good properties derived for probabilities; we finish the section showing that these conditions are nevertheless necessary. In Section 4 we treat the problem of obtaining the set of dominating \( k \)-additive belief functions dominating a capacity. Section 5 is devoted to the applications in decision making. We finish with the conclusions and open problems.

## 2 Basic concepts on non-additive measures

Consider a finite referential set of \( n \) elements (criteria in Multicriteria Decision Making, players in Cooperative Games, incomes in Welfare Theory, ...), \( X = \{1, ..., n\} \). The set of subsets of \( X \) is denoted by \( \mathcal{P}(X) \), while the set of subsets whose cardinality is less or equal than \( k \) is denoted by \( \mathcal{P}^{k}(X) \). Subsets of \( X \) are denoted \( A, B, ... \). We will sometimes write \( i_{1} \cdots i_{k} \) instead of \( \{i_{1}, ..., i_{k}\} \) in order to avoid heavy notation; brackets are usually omitted for singletons and subsets of two elements.

In order to be self-contained, we introduce the concepts that will be needed throughout the paper.

**Definition 1.** A non-additive measure \([4]\) or fuzzy measure \([27]\) or capacity \([2]\) over \( (X, \mathcal{P}(X)) \) is a mapping \( \mu : \mathcal{P}(X) \to [0,1] \) satisfying

- \( \mu(\emptyset) = 0, \mu(X) = 1 \) (boundary conditions).
- \( \forall A, B \in \mathcal{P}(X), \text{ if } A \subseteq B, \text{ then } \mu(A) \leq \mu(B) \) (monotonicity).

In the sequel, we will mainly use the word ”capacity”. We denote the set of all capacities over \( X \) by \( \mathcal{FM}(X) \).

An important class of capacities is the following:

**Definition 2.** A unanimity game over \( A \subseteq X, A \neq \emptyset \) is a capacity defined by

\[
\nu_{A}(B) := \begin{cases} 
1 & \text{if } A \subseteq B \\
0 & \text{otherwise}
\end{cases}
\]

For \( \emptyset \), we define the unanimity game by

\[
\nu_{\emptyset}(B) := \begin{cases} 
1 & \text{if } B \neq \emptyset \\
0 & \text{if } B = \emptyset
\end{cases}
\]
Definition 3. [5] Let \( \mu \) be a capacity. We say that \( \mu \) is \( k \)-monotone (for \( k \geq 2 \)) if and only if for any family of \( k \) subsets \( A_1, \ldots, A_k \) of \( X \),
\[
\mu \left( \bigcup_{j=1}^k A_j \right) \geq \sum_{\emptyset \not\subseteq I \subseteq \{1, \ldots, k\}} (-1)^{|I|+1} \mu \left( \bigcap_{j \in I} A_j \right).
\]
If \( \mu \) is \( k \)-monotone for all \( k \geq 2 \), then it is said that \( \mu \) is a belief function.

We will denote by \( \mathcal{BE}(X) \) the set of all belief functions on \( X \).

There are other set functions that can be used to equivalently represent a capacity. In this paper we will need the so-called Möbius transform.

Definition 4. [24] Let \( \mu \) be a set function (not necessarily a non-additive measure) on \( X \). The Möbius transform (or inverse) of \( \mu \) is another set function on \( X \) defined by
\[
m(A) := \sum_{B \subseteq A} (-1)^{|A\setminus B|} \mu(B), \quad \forall A \subseteq X.
\]
The Möbius transform given, the original set function can be recovered through the Zeta transform [1]:
\[
\mu(A) = \sum_{B \subseteq A} m(B).
\]
The value \( m(A) \) represents the strength of the subset \( A \) in any coalition in which it appears. The Möbius transform corresponds to the basic probability mass assignment in Dempster-Shafer theory of evidence [25].

The Möbius transform can be applied to any set function. Given a Möbius inverse, necessary and sufficient conditions for obtaining a proper capacity have been investigated. The following can be proved:

**Proposition 1.** [1] A set of \( 2^n \) coefficients \( m(A), A \subseteq X \), corresponds to the Möbius transform of a capacity if and only if

(i) \( m(\emptyset) = 0 \), \( \sum_{A \subseteq X} m(A) = 1 \).

(ii) \( \sum_{i \in B \subseteq A} m(B) \geq 0 \), for all \( A \subseteq X \), for all \( i \in A \).

In particular, (ii) implies that \( m(i) \geq 0 \), \( \forall i \in X \). Related to unanimity games and belief functions, the following result can be proved:

**Proposition 2.** [12, 25]

- The Möbius transform of the unanimity game on \( A \), with \( A \neq \emptyset \), is given by
  \[
  m(B) = \begin{cases} 
  1 & \text{if } B = A \\
  0 & \text{otherwise}
  \end{cases}
  \]
  For \( u_\emptyset \), the Möbius inverse is given by \( m(A) = (-1)^{|A|+1} \), \( \forall A \subseteq X, A \neq \emptyset \) and \( m(\emptyset) = 0 \).

- \( \mu \) is a belief function if and only if its corresponding Möbius inverse is non-negative.

In order to determine a capacity, \( 2^n - 2 \) values are necessary. The number of coefficients grows exponentially with \( n \) and so does the complexity of the problem of identification. This drawback reduces considerably the practical use of non-additive measures. In an attempt to reduce complexity, some subfamilies of non-additive measures have been defined e.g. \( k \)-additive measures [9], \( p \)-symmetric measures [18], \( k \)-intolerant measures [15], \( \lambda \)-measures [28], or more generally decomposable measures [6]. In this paper we will consider \( k \)-additive measures.

**Definition 5.** [9] A non-additive measure \( \mu \) is said to be \( k \)-order additive or \( k \)-additive if its Möbius transform vanishes for any \( A \subseteq X \) such that \( |A| > k \) and there exists at least one subset \( A \) with exactly \( k \) elements such that \( m(A) \neq 0 \).
In this sense, a probability measure is just a 1-additive measure [9]. Thus, k-additive measures generalize probability measures, that are very restrictive in many situations as they do not allow interactions between the elements of $X$. They fill the gap between probability measures and general non-additive measures. For a k-additive measure, the number of coefficients is reduced to

$$\sum_{i=1}^{k} \binom{n}{i}.$$ 

More about k-additive measures can be found e.g. in [10]. We will denote the set of all k'-additive measures with $k' \leq k$ on $X$ by $\mathcal{FM}_k^{k'}(X)$; in particular, $\mathcal{FM}_1^{1}(X)$ represents the set of all probability distributions. Analogously, the set of all k'-additive belief functions, with $k' \leq k$ on $X$ is denoted by $\mathcal{BEL}_k^{k'}(X)$. Remark that $\mathcal{FM}_1^{1}(X) = \mathcal{BEL}_1^{1}(X)$.

**Definition 6.** Consider a capacity $\mu$. We say that a capacity $\mu^*$ **dominates** $\mu$, and we denote it $\mu^* \geq \mu$, if and only if

$$\mu^*(A) \geq \mu(A), \forall A \subseteq X.$$ 

Given a capacity $\mu$, the set of all k-additive (at most) capacities dominating $\mu$ will be denoted by $\mathcal{FM}_k^k(\mu)$. The set of all k-additive (at most) belief functions dominating $\mu$ will be denoted by $\mathcal{BEL}_k^k(\mu)$. It is trivial to see from Definition 6 that $\mathcal{FM}_k^k(\mu)$ (resp. $\mathcal{BEL}_k^k(\mu)$) is a convex polyhedron and, consequently, it is characterized by its vertices. The set of vertices is called the **profile**.

Let $\Sigma$ be the set of all permutations on $X$; given $\sigma \in \Sigma$, let us denote by $X^l_\sigma$, $0 \leq l \leq n$, the subset of $X$ defined by

$$X^l_\sigma := \{\sigma(1), ..., \sigma(l)\}, l \geq 1, X^0_\sigma = \emptyset.$$ 

**Definition 7.** Given a capacity $\mu$, we denote by $\mathcal{FM}_k^k(\mu)$ (resp. $\mathcal{BEL}_k^k(\mu)$) the subset of $\mathcal{FM}_k^k(X)$ (resp. $\mathcal{BEL}_k^k(X)$) containing the capacities (resp. belief functions) $\mu^*$ satisfying

$$\exists \sigma \in \Sigma | \forall l = 1, ..., n, \mu^*(X^l_\sigma) = \mu(X^l_\sigma).$$

### 3 Dominating probabilities versus dominating k-additive measures

Let us now start our study. Chateauneuf and Jaffray have studied in [1] the problem of dominating a capacity by probabilities. They obtain the following result:

**Theorem 1.** Let $\mu$ be a capacity on $X$, $m$ its Möbius transform, and suppose $P \in \mathcal{FM}_1^1(\mu)$. Then, $P$ can be put under the following form:

$$P(\{i\}) = \sum_{B \ni i} \lambda(B, i)m(B), \forall i \in X,$$

and $P(A) = \sum_{i \in A} P(\{i\})$ for any $A \subseteq X$. The function $\lambda : \mathcal{P}(X) \times X \rightarrow [0, 1]$ is a weight function satisfying:

$$\sum_{i \in B} \lambda(B, i) = 1, \forall B \subseteq X.$$

$$\lambda(B, i) = 0 \text{ whenever } i \notin B.$$ 

Dempster has shown the same result in [3] and also Shapley in [26], but both of them only for belief functions. Grabisch has extended Theorem 1 for the k-additive case. He proved the following result:

**Theorem 2.** [13] Let $\mu$ be a capacity on $X$, $m$ its Möbius transform and suppose that $\mu^* \in \mathcal{FM}_k^k(\mu)$. Then, necessarily, the Möbius transform $m^*$ of $\mu^*$ can be put under the following form:

$$m^*(A) = \sum_{B \cap A \neq \emptyset} \lambda(B, A)m(B), \forall A \in \mathcal{P}^k(X),$$
where function \( \lambda : \mathcal{P}(X) \times \mathcal{P}(X) \rightarrow \mathbb{R} \) is such that
\[
\sum_{A \mid B \cap A \neq \emptyset} \lambda(B, A) = 1, \forall B \subseteq X. \tag{1}
\]
\[
\lambda(B, A) = 0, \forall A \in \mathcal{P}(X), A \cap B = \emptyset. \tag{2}
\]

In the following, the set function obtained for a given choice of \( \lambda \) in the conditions of Theorem 2 will be denoted by \( \mu_\lambda \) and \( m_\lambda \) its corresponding Möbius transform.

At first sight, Theorem 2 is a straight generalization of Theorem 1 and both results are very similar. However, there are two important facts related to \( \lambda \) that should be kept in mind:

- \( \lambda \) in Theorem 2 can attain negative values, while in Theorem 1 it is a weight function.
- The condition \( B \ni i \) of Theorem 1 can be generalized in two different ways for general subsets: \( A \subseteq B \) and \( A \cap B \neq \emptyset \). The second case includes the first one. It must be remarked that in Theorem 2 the second alternative is used.

These two conditions are necessary for the general \( k \)-additive case and cannot be simplified, as the following examples show:

**Example 1.** [13] Let us start showing that \( A \subseteq B \) instead of \( A \cap B \neq \emptyset \) in Equations (1) and (2) does not suffice to recover all \( k \)-additive measures dominating \( \mu \). Consider \( |X| = 3 \) and the fuzzy measure defined by:

\[
\begin{array}{cccccccc}
 & \{1\} & \{2\} & \{3\} & \{1, 2\} & \{1, 3\} & \{2, 3\} & \{1, 2, 3\} \\
m & 0.1 & 0.1 & 0.1 & 0.2 & 0.2 & 0.2 & 0.1 \\
\mu & 0.1 & 0.1 & 0.1 & 0.3 & 0.4 & 0.4 & 1
\end{array}
\]

Let us consider now the belief function defined by:

\[
\begin{array}{cccccccc}
 & \{1\} & \{2\} & \{3\} & \{1, 2\} & \{1, 3\} & \{2, 3\} & \{1, 2, 3\} \\
m^* & 0.35 & 0.1 & 0.1 & 0 & 0 & 0.45 & 0 \\
\mu^* & 0.35 & 0.1 & 0.1 & 0.45 & 0.45 & 0.65 & 1
\end{array}
\]

Then, \( \mu^* \geq \mu \) and \( \mu^* \) is 2-additive, but in the sharing procedure, if we change \( A \cap B \) by \( A \subseteq B \), \( m^*(2, 3) \) should be derived only from \( m(2, 3) \) and \( m(1, 2, 3) \). But then, \( m(2, 3) + m(1, 2, 3) = 0.4 < m^*(2, 3) = 0.45 \), and therefore no sharing function would lead to \( \mu^* \).

**Example 2.** Let us now show that \( \lambda \) may attain negative values. Consider \( |X| = 3 \) and again the capacity \( \mu \) defined in Example 1. Let us take now the 2-additive measure defined by:

\[
\begin{array}{cccccccc}
 & \{1\} & \{2\} & \{3\} & \{1, 2\} & \{1, 3\} & \{2, 3\} & \{1, 2, 3\} \\
m^* & 0.7 & 0.3 & 0.2 & -0.2 & 0 & 0 & 0 \\
\mu^* & 0.7 & 0.3 & 0.2 & 0.8 & 0.9 & 0.5 & 1
\end{array}
\]

Thus, \( \mu^* \geq \mu \) but \( m^*(1, 2) < 0 \) and then we need a function \( \lambda \) taking negative values.

The following notations will be useful to shed light on the differences between these two theorems:

\[
\Lambda_{k}^V := \{ \lambda : \mathcal{P}(X) \times \mathcal{P}(X) \rightarrow \mathbb{R} \mid \forall B \in \mathcal{P}(X), \sum_{A \cap B \neq \emptyset} \lambda(B, A) = 1, \lambda(B, A) = 0 \text{ if } A \cap B = \emptyset \}.
\]

\[
\Lambda_{k,+}^V := \{ \lambda : \mathcal{P}(X) \times \mathcal{P}(X) \rightarrow [0, 1] \mid \forall B \in \mathcal{P}(X), \sum_{A \cap B \neq \emptyset} \lambda(B, A) = 1, \lambda(B, A) = 0 \text{ if } A \cap B = \emptyset \}.
\]

\[
\Lambda_{k,\leq,+}^V := \{ \lambda : \mathcal{P}(X) \times \mathcal{P}(X) \rightarrow [0, 1] \mid \forall B \in \mathcal{P}(X), \sum_{A \cap B \neq \emptyset} \lambda(B, A) = 1, \lambda(B, A) = 0 \text{ if } A \subseteq B \}.
\]
Given \( \mu \in \mathcal{FM}(X) \) whose corresponding Möbius transform is \( m \), we define
\[
\mathcal{M}_{\Lambda_k}^\mu := \{ \mu_\lambda \mid m_\lambda(A) = \sum_{B \in \mathcal{P}(X)} \lambda(B, A)m(B), \lambda \in \Lambda_k \}. 
\]
\[
\mathcal{M}_{\Lambda_k,+}^\mu := \{ \mu_\lambda \mid m_\lambda(A) = \sum_{B \in \mathcal{P}(X)} \lambda(B, A)m(B), \lambda \in \Lambda_k^{+,+} \}. 
\]
\[
\mathcal{M}_{\Lambda_{\preceq,+}}^\mu := \{ \mu_\lambda \mid m_\lambda(A) = \sum_{B \in \mathcal{P}(X)} \lambda(B, A)m(B), \lambda \in \Lambda_k_{\preceq,+} \}. 
\]
With this notation, \( \mathcal{M}_{\Lambda_{\preceq,+}}^\mu \) is the set of functions in the conditions of Theorem 1 and \( \mathcal{M}_{\Lambda_k}^\mu \) represents the set of functions obtained by Theorem 2.

For these families, Grabisch has studied in [11] their algebraic structure. Let us define a “composition” \( \star \) between two functions \( \lambda, \lambda' \) as
\[
\lambda \star \lambda'(A, B) = \sum_{C \in \mathcal{P}(X)} \lambda(A, C)\lambda'(C, B), \forall (A, B) \in \mathcal{P}(X) \times \mathcal{P}^k(X). 
\]
Then, the following hold:

**Proposition 3. [11]**
- \((\Lambda_k^{+,+}, \star)\) and \((\Lambda_k^{+,+}, \star)\) are not stable under \( \star \).
- \((\Lambda_k^+, \star)\) is a monoid.
- \((\Lambda_k^{+,+}, \star)\) is a monoid.

It can be shown [1] that in general \( \mathcal{FM}_{\preceq}^k(\mu) \subseteq \mathcal{M}_{\Lambda_{\preceq,+}}^\mu \) and equality cannot be ensured for any measure \( \mu \); the same can be easily seen for \( \mathcal{M}_{\Lambda_k}^\mu \) and \( \mathcal{FM}_{\preceq}^k(\mu) \). As we will see later, \( \mathcal{M}_{\Lambda_k}^\mu \) provides a very wide set of functions and therefore we obtain some unpleasant results. Then, other questions arise:

**Question 1.** Are the sets of functions obtained in Theorems 1 and 2 included in the set \( \mathcal{FM}_{\preceq}^k(X) \), i.e. under which conditions on \( \mu \) can we ensure that all functions obtained in Theorem 2 are \( k \)-additive capacities?

For probabilities, the following can be proved:

**Proposition 4. [1]** Let \( \mu \) be a capacity; \( m \) its Möbius inverse. Then, \( \mathcal{M}_{\Lambda_{\preceq,+}}^\mu \subseteq \mathcal{FM}_{\preceq}^k(X) \) if and only if
\[
m(i) + \sum_{B \ni i, B \neq \{i\}} \min\{m(B), 0\} \geq 0, \forall i \in X. 
\]
A similar result does not hold for the \( k \)-additive case, as next proposition shows:

**Proposition 5.** Let \( \mu \) be a capacity. Then, for \( k \geq 2 \), we can always find \( \lambda \in \Lambda_k^+ \) such that \( \mu_\lambda \notin \mathcal{FM}_{\preceq}^k(X) \).

**Proof:** Let \( m \) be the Möbius inverse of \( \mu \). Then, by Theorem 2 we have for each \( \lambda \in \Lambda_k^+ \),
\[
m_\lambda(i) \geq 0 \iff \sum_{B \ni i \in B} \lambda(B, i)m(B) \geq 0. 
\]
As \( \sum_{B \in \mathcal{X}} m(B) = 1 \) by Proposition 1, we can always find a subset \( A \) such that \( m(A) > 0 \). Consider \( i \in A \) and \( j \neq i \).

We make the following choice of \( \lambda \):
\[
\lambda(A, i) = -1, \lambda(A, ij) = 2, \lambda(A, C) = 0, \forall C \neq \{i\}, \{i, j\}. 
\]
Now, for any subset $B$ such that $B \neq A$, $B \ni i$, define

$$\lambda(B, i) = 0, \lambda(B, ij) = 1, \lambda(B, C) = 0, \forall C \neq \{i\}, \{i, j\}.$$  

Finally, for $B \subseteq X$ s.t. $i \notin B$, we choose $k \in B$ and define

$$\lambda(B, k) = 1, \lambda(B, C) = 0, \forall C \neq \{k\}.$$  

It is clear that $\lambda$ satisfies the conditions of Theorem 2. But for this choice of $\lambda$, we have

$$m_\lambda(i) = \sum_{i \in B} \lambda(B, i)m(B) = -m(A) < 0,$$

and thus, $\mu_\lambda$ is not a capacity by Proposition 1.

We have seen that the monotonicity conditions fail. Nevertheless, other conditions hold for any $\mu$ and all possible choices of $\lambda \in \Lambda_k^\pm$.

**Lemma 1.** Let $\mu^* \in M_{\Lambda_k}(\mu)$ and let $m^*$ be its corresponding Möbius inverse. Then,

1. $m^*(\emptyset) = 0$.
2. $\sum_{A \subseteq X} m^*(A) = 1$.

**Proof:** 1 is trivial. For 2, it suffices to consider

$$\sum_{A \subseteq X} m^*(A) = \sum_{A \subseteq X} \sum_{B \mid B \cap A \neq \emptyset} \lambda(B, A)m(B) = \sum_{B \subseteq X} m(B) \sum_{A \mid B \cap A \neq \emptyset} \lambda(B, A) = \sum_{B \subseteq X} m(B) = 1,$$

whence the result.

**Question 2.** What is the profile of $\mathcal{FM}_\geq^k(\mu)$?

The corresponding result for probabilities is

**Proposition 6.** \cite{1} If $\mu$ is a 2-monotone capacity, then $\mathcal{FM}_\geq^1(\mu)$ is the profile of $\mathcal{FM}_\geq^1(\mu)$, i.e. $\mathcal{FM}_\geq^1(\mu)$ is the convex closure of $\mathcal{FM}_\geq^1(\mu)$.

We cannot translate Proposition 6 for the general $k$-additive case as next proposition shows. We start with a preliminary lemma.

**Lemma 2.** For all $\mu \in \mathcal{FM}(X)$, $\exists A \subseteq X$ such that $m(A) > 0$ and $m(B) = 0$, $\forall B \subset A$.

**Proof:** Let $\mu$ be a capacity and $m$ its corresponding Möbius inverse. Let us start proving that we can always find $A \subseteq X$ such that $m(A) \neq 0$ and $m(B) = 0$, $\forall B \subset A$.

If there exists $i \in X$ such that $m(i) \neq 0$, the result is proved.

Otherwise, $m(i) = 0$, $\forall i \in X$. Let us then consider $m(i, j)$, $\forall i, j \in X$. If there exists $\{i, j\} \subseteq X$ such that $m(i, j) \neq 0$ the result holds.

Taking into account that $\sum_{A \subseteq X} m(A) = 1$ (Proposition 1), if we repeat this process, we obtain that there exists a subset $A$ such that $m(A) \neq 0$ and $m(B) = 0$, $\forall B \subset A$.

Now, it suffices to show that $m(A) > 0$. By the monotonicity constraints on $m$ (Proposition 1), for $i \in A$ we have

$$\sum_{i \in B \subseteq A} m(B) \geq 0.$$

But, $\sum_{i \in B \subseteq A} m(B) = m(A)$, whence the result.
Proposition 7. Let \( \mu \) be a capacity, \( \mu \neq u_X \). Then, for \( k \geq 2 \), there exists a \( k \)-additive (at most) capacity \( \mu^* \) in \( \mathcal{F}\mathcal{M}^k_\geq(\mu) \) such that \( \mu^* \notin \mathcal{F}\mathcal{M}^k_\geq(\mu) \).

**Proof:** By Lemma 2, we can always find a subset \( A \subseteq X \) such that \( m(A) > 0 \) and \( m(B) = 0, \forall B \subset A \). Let us suppose without loss of generality \( A = \{1, \ldots, l\} \). Then, consider the permutation \( \sigma \) defined by
\[
\sigma(1) = l + 1, \sigma(2) = l + 2, \ldots, \sigma(n - l) = n, \sigma(n - l + 1) = l, \sigma(n - l + 2) = l - 1, \ldots, \sigma(n) = l.
\]
Define a function \( \mu^* \) by the following Möbius transform:
\[
m^*(l + 1) = \mu(l + 1).
\]
\[
m^*(i) = \mu(l + 1, \ldots, i) - \mu(l + 1, \ldots, i - 1), l + 2 \leq i \leq n.
\]
\[
m^*(i, n) = \mu(i, \ldots, n) - \mu(i + 1, \ldots, n), l \geq i \geq 1,
\]
and take \( m^*(C) = 0 \) otherwise.

As \( \mu \) is a capacity, \( m^*(C) \geq 0, \forall C \subseteq X \). On the other hand,
\[
\sum_{A \subseteq X} m^*(A) = \sum_{i=l+1}^n m^*(i) + \sum_{i=1}^l m^*(i, n) = \mu(l + 1, \ldots, n) + \mu(1, \ldots, n) - \mu(l + 1, \ldots, n) = 1.
\]
Thus, \( \mu^* \) is a capacity (it is indeed a belief function). Moreover, \( \mu^* \in \mathcal{F}\mathcal{M}^2(X) \). Finally, \( \mu^* \in \mathcal{F}\mathcal{M}^k_\geq(\mu) \), as it takes the same values of \( \mu \) for \( \mathcal{X}_k \), \( \forall j = 1, \ldots, n \).

However, \( \mu^* \notin \mathcal{F}\mathcal{M}^k_\geq(\mu) \) because
\[
\mu^*(A) = \sum_{B \subseteq A} m^*(B) = 0 < m(A) = \sum_{B \subseteq A} m(B) = \mu(A),
\]
whence the result.

Question 3. When do all functions obtained in Theorems 1 and 2 dominate \( \mu \)?

The corresponding result for probabilities is

Proposition 8. [1] Let \( \mu \) be a capacity. Then \( \mathcal{M}_{\Lambda^k_{\geq+}}(\mu) = \mathcal{F}\mathcal{M}^1_\geq(\mu) \) if and only if \( \mu \) is a belief function.

We have shown in Proposition 5 that we can always find \( \lambda \in \Lambda^k_{\geq} \) such that \( \mu_{\lambda} \) is not a capacity. Thus, we cannot obtain a result similar to the one in Proposition 8 for Theorem 2. However, we can ask ourselves what happens if we restrict to capacities in the conditions of Theorem 2, i.e. if \( \mu^* \in \mathcal{M}_{\Lambda^k_{\geq}}(\mu) \cap \mathcal{F}\mathcal{M}^k(X) \). Even in that case, we can find \( \lambda \in \Lambda^k_{\geq} \) such that \( \mu_{\lambda} \) is a capacity but it does not dominate \( \mu \) whenever \( \mu \neq u_X \) (remark that \( u_X \) is dominated by all other capacities):

Proposition 9. Let \( \mu \) be a capacity, \( \mu \neq u_X \). Then, for \( k \geq 2 \), there exists \( \lambda \in \Lambda^k_{\geq} \) such that \( \mu_{\lambda} \in \mathcal{F}\mathcal{M}^k(X) \) but \( \mu_{\lambda} \notin \mathcal{F}\mathcal{M}^k_\geq(\mu) \).

**Proof:** Consider \( \sigma \) and \( A \) defined as in Proposition 7. Consider also the capacity \( \mu^* \) defined in that proposition. We have shown in Proposition 7 that \( \mu^* \) is a 2-additive belief function and that \( \mu^* \notin \mathcal{F}\mathcal{M}^k_\geq(\mu) \). Then, it suffices to find a choice of \( \lambda \in \Lambda^k_{\geq} \) such that \( \mu_{\lambda} = \mu^* \).

For each \( B = \{\sigma(i_1), \ldots, \sigma(i_r)\} \subseteq X \), with \( i_1 < i_2 < \ldots < i_r \), we denote \( i_B = \sigma(i_r) \), i.e. the element such that \( \sigma(\max i) = i_B \). Then, let us define \( \lambda \) as follows:
\[
\lambda(B, i_B) = 1, \lambda(B, C) = 0, \text{ for } C \neq \{i_B\} \text{ if } l + 1 \leq i_B \leq n.
\]
\[
\lambda(B, \{i_B, n\}) = 1, \lambda(B, C) = 0, \text{ otherwise if } 1 \leq i_B \leq l.
\]
It is clear that \( \lambda \in \Lambda^k_{\geq} \) and it is straightforward to see that \( \mu_{\lambda} = \mu^* \).

To end this section, let us consider the following curious result:
Example 3. Consider $|X| = 2$ and consider the capacity $u_\emptyset$. The Möbius inverse $m$ of $\mu$ is given by $m(1) = 1$, $m(2) = 1$, $m(1,2) = -1$ (by Proposition 2). It is easy to see that all capacities are in the conditions of Theorem 2. However, the only capacity dominating $\mu$ is $\mu$ itself.

We have seen that we cannot derive similar results for the $k$-additive case as those obtained for probabilities. The reason relies on the fact that $\mathcal{M}_{\Lambda_+^k}(\mu)$ is very large for $k \geq 2$; however, Examples 1 and 2 show that Equations (1) and (2) cannot be strengthened. Consequently, we have to look for other results in order to obtain conditions for $\lambda$ more similar to those of Theorem 1. This leads us to add other constraints on the dominating measure.

4 Characterization of dominating belief functions

In the previous section we have seen that it is not possible to derive similar properties for Theorem 2 to those obtained for Theorem 1. The reason relies in the fact that $\Lambda_+^k$, although generalizing $\Lambda_+^1$, provides a very wide class of functions. In other words, the ideal generalization for $\Lambda_+^k$ should be $\Lambda_+^{k+1}$. However, we have already seen (Examples 1 and 2) that this set does not cope in general with the set of all dominating $k$-additive measures, i.e. the conditions on $\lambda$ given in Theorem 2 cannot be strengthened. This means that $k$-additive measures constitute a too wide generalization of probabilities. Then, in order to derive a similar result to Theorem 1 keeping these good properties, we have to impose some other constraints on the dominating measures.

In this sense, note that probabilities can be seen not only as 1-additive measures but also as 1-additive belief functions. Then, we are going to restrict ourselves to the case of finding the set of $k$-additive belief functions dominating a given capacity $\mu$. In this case, our next result proves that we can take positive weight functions:

Theorem 3. Let $\mu, m, \mu^* : \mathcal{P}(X) \to \mathbb{R}$, where $\mu$ is a capacity, $m$ its Möbius inverse, and $\mu^* \in \mathcal{BE}\ell_\geq^k(\mu)$. Then, necessarily the Möbius transform $m^*$ of $\mu^*$ can be put under the following form:

$$m^*(A) = \sum_{B : A \cap B \neq \emptyset} \lambda(B,A)m(B), \forall A \in \mathcal{P}^k(X),$$

where function $\lambda : \mathcal{P}(X) \times \mathcal{P}^k(X) \to [0,1]$ is such that

$$\sum_{A : B \cap A \neq \emptyset} \lambda(B,A) = 1, \forall B \in \mathcal{P}(X).$$

(3)

$$\lambda(B,A) = 0, \text{ if } A \cap B = \emptyset.$$  

(4)

Proof: We will follow the same sketch of proof as Chateauneuf and Jaffray in [1] and Grabisch in [13] for Theorems 1 and 2, respectively.

Consider then the transhipment problem on a capacitated network consisting of:

- A set of sources $E = \{e_B : B \in \mathcal{P}(X), m(B) > 0\}$ with supply $m(B)$ at $e_B \in E$.
- A set of sinks $E^* \cup S$ where $E^* = \{e_B^* : B \in \mathcal{P}^k(X)\}$ with demand $m^*(B)$ at $e_B^* \in E^*$, and $S = \{s_B : B \in \mathcal{P}(X), m(B) < 0\}$ with demand $-m(B)$ at $s_B \in S$.
- Arcs, with infinite capacity, joining a source $e_B \in E$ to a sink $e_C^* \in E^*$ or a sink $e_C^* \in E^*$ to a sink $s_B \in S$. Then, we have $B \cap C \neq \emptyset$.

Figure 1 shows an example of such a flow network. Note that there is no excess supply, since

$$\sum_{A \in \mathcal{P}^k(X)} m^*(A) = 1 = \sum_{A \in \mathcal{P}(X)} m(A) = \sum_{e_B \in E} m(B) - \sum_{s_B \in S} (-m(B)).$$

Thus, a feasible flow $\phi$, has to saturate the supply and demand constraints:

$$\sum_{A : A \cap B \neq \emptyset, e_A^* \in E^*} \phi(e_B, e_A^*) = m(B), \forall e_B \in E.$$
To any feasible flow $\phi$, one can associate a function $\lambda$ partially defined by:

$$
\lambda(B, A) = \begin{cases} 
\frac{\phi(e_B, e_A^*)}{m(B)} & \text{for } m(B) > 0 \\
\frac{\phi(e_A^*, s_B)}{-m(B)} & \text{for } m(B) < 0 
\end{cases}
$$

choosing, for $B$ such that $m(B) = 0$, arbitrary $\lambda(B, A)$ satisfying Equations (3) and (4). Thus, all we need to prove is the existence of such a feasible flow in the network. According to Gale’s theorem [8], this amounts to check that, for each partition $\{\bar{N}, N\}$ of the set of nodes, $k(\bar{N}, N) \geq d(N)$, where $k(\bar{N}, N)$ is the sum of capacities of the arcs joining a node in $\bar{N}$ to a node in $N$, and $d(N)$ is the net demand in $N$, (i.e. the difference between the sum of the demands and the sum of the supplies at the various nodes in $N$).

This inequality is obviously satisfied when $k(\bar{N}, N) = \infty$, i.e. when there exists $e_B \in E \cap \bar{N}$ and $e_A^* \in E^* \cap N$, $A \cap B \neq \emptyset$, or there exists $e_{\bar{B}}^* \in E^* \cap \bar{N}$ and $s_B \in S \cap \bar{N}$, $A \cap B \neq \emptyset$.

In the alternative case, where $k(\bar{N}, N) = 0$, it can be noted:

$$
d(N) = \sum_{e_A^* \in E^* \cap \bar{N}} m^*(A) = \sum_{e_A^* \in E^* \cap N} m(A) + \sum_{s_B \in S \cap N} (-m(B)).
$$

Consider $H^* := \bigcup_{e_A^* \in E^* \cap N} A$. Then, as $\mu^*$ is a belief function $\sum_{A \subseteq H^*} m^*(A) \geq 0$, and thus,

$$
\sum_{e_A^* \in E^* \cap N} m^*(A) \leq \mu^*(H^*).
$$

On the other hand, again as $\mu^*$ is a belief function, $\mu^*(H^*) + \mu^*(\bar{H}^*) \leq 1$, and thus $\mu^*(H^*) \leq 1 - \mu^*(\bar{H}^*)$, whence

$$
\sum_{e_A^* \in E^* \cap N} m^*(A) \leq 1 - \mu^*(\bar{H}^*).
$$

Now,

$$
\sum_{e_A \in E \cap N} m(A) + \sum_{s_A \in S \cap N} m(A) = 1 - \sum_{e_A \in E \cap \bar{N}} m(A) - \sum_{s_A \in S \cap \bar{N}} m(A).
$$
Let us show that $e_B \in \mathcal{E} \cap \overline{\mathcal{N}}$ implies $B \cap \mathcal{H}^* = \emptyset$. Otherwise, $B \cap \mathcal{H}^* \neq \emptyset$, and then $\exists i \in B \cap \mathcal{H}^*$. On the other hand, by definition of $\mathcal{H}^*$, $\exists A$ such that $e_A^* \in \mathcal{E} \cap \overline{\mathcal{N}}$ satisfying $i \in A$. This implies the existence of an arc joining $e_B$ and $e_A^*$ and $k(\overline{\mathcal{N}}, \mathcal{N}) = \infty$, contradicting $k(\overline{\mathcal{N}}, \mathcal{N}) = 0$.

We deduce then that $B \subseteq \mathcal{H}^*$ and consequently,

$$\sum_{e_B \in \mathcal{E} \cap \overline{\mathcal{N}}} m(B) \leq \sum_{B \subseteq \mathcal{H}^*, m(B) > 0} m(B).$$

Remark on the other hand that if $B \subseteq \mathcal{H}^*$ and $m(B) < 0$, then $\forall A \subseteq X$ such that $A \cap B \neq \emptyset$, necessarily $e_A^* \in \overline{\mathcal{N}}$. Otherwise, $\exists A \subseteq X$ such that $A \cap B \neq \emptyset$ and $e_A^* \in \mathcal{N}$. But then, $A \subseteq \mathcal{H}^*$, and therefore $B \cap \mathcal{H}^* \neq \emptyset$, contradicting $B \subseteq \mathcal{H}^*$.

We conclude that for such a $B$, we have $s_B \in S \cap \overline{\mathcal{N}}$. Otherwise $s_B \in S \cap \mathcal{N}$, and as there exists an arc joining $e_A^*$ and $s_B$ for $A$ such that $A \cap B \neq \emptyset$, we would obtain $k(N, \overline{\mathcal{N}}) = \infty$, contradicting $k(N, \overline{\mathcal{N}}) = 0$. Then,

$$\sum_{A \cap s_A \in S \cap \overline{\mathcal{N}}} m(A) \leq \sum_{A \subseteq \mathcal{H}^*, m(A) < 0} m(A).$$

As a consequence,

$$1 - \sum_{e_A \in \mathcal{E} \cap \overline{\mathcal{N}}} m(A) - \sum_{s_A \in S \cap \overline{\mathcal{N}}} m(A) \geq 1 - \sum_{A \subseteq \mathcal{H}^*, m(A) > 0} m(A) - \sum_{A \subseteq \mathcal{H}^*, m(A) < 0} m(A) = 1 - \sum_{A \subseteq \mathcal{H}^*} m(A) = 1 - \mu(\mathcal{H}^*).$$

Finally,

$$d(\mathcal{N}) \leq 1 - \mu^*(\mathcal{H}^*) - 1 + \mu(\mathcal{H}^*) = -\mu^*(\mathcal{H}^*) + \mu(\mathcal{H}^*) \leq 0,$$

as $\mu^* \geq \mu$.

Following the notation introduced in the previous section, the set of functions obtained in Theorem 3 is $\mathcal{M}_{\Lambda_{\mathcal{H}^*}^k}(\mu)$.

We have to remark that the condition $A \cap B \neq \emptyset$ is necessary and cannot be replaced by $A \subseteq B$, even in the special case of $\mu$ being a belief function, as the following example shows:

**Example 4.** Consider $|X| = 4$ and the 2-additive belief function $\mu$ whose Möbius inverse is given by:

$$m(i) = 0.1, m(i, j) = 0.1, \forall i, j \in X.$$

Consider now the belief function $\mu^*$ whose Möbius transform is given by

$$m^*(i) = 0.2, \forall i \in X, m^*(1, 2, 3) = 0.2.$$

Then, $\mu^*$ is a 3-additive belief function dominating $\mu$. However, as $m(A) = 0, \forall A \setminus \{1, 2, 3\} \subseteq A$, we have to consider $\Lambda_{\mathcal{H}^*}^k$ in order to derive a function $\lambda$ such that $\lambda \subseteq \mu^*$. As $\Lambda_{\mathcal{H}^*}^k \subseteq \Lambda_{\mathcal{H}^*}^k$, we have $\mathcal{M}_{\Lambda_{\mathcal{H}^*}^k}(\mu) \subseteq \mathcal{M}_{\Lambda_{\mathcal{H}^*}^k}(\mu)$. Consequently, Theorem 3 provides a better result than Theorem 2 for belief functions. However, $\mathcal{M}_{\Lambda_{\mathcal{H}^*}^k}(\mu)$ is still too large, as it is shown in next proposition. We need a preliminary lemma.

**Lemma 3.** If $\mu$ is a belief function, then all measures in $\mathcal{M}_{\Lambda_{\mathcal{H}^*}^k}(\mu)$ are belief functions.

**Proof:** As $\mu$ is a belief function, $m(A) \geq 0, \forall A \in \mathcal{P}(X)$ by Proposition 2. Moreover, $\lambda(B, A) \geq 0, \forall A, B$ and consequently, $m^*(A) \geq 0, \forall A \in \mathcal{P}^k(X)$.

Thus, it suffices to show that $\sum_{B \subseteq X} m^*(B) = 1$, but this is true by Lemma 1.

**Proposition 10.** Let $\mu$ be a capacity. Then, $\mathcal{M}_{\Lambda_{\mathcal{H}^*}^k}(\mu) = \mathcal{B} \mathcal{L}_{\geq}^k(\mu)$ if and only if $\mu = u_X$. 

11
Proof: If \( \mu = u_X \in \mathcal{BEL}(X) \), then \( \mathcal{M}_{\Lambda_{\mathbb{N}^+}^k}(u_X) \subseteq \mathcal{BE}_{\geq k}(X) \) by Lemma 3. Moreover, as \( u_X \) is dominated by any other capacity, we conclude \( \mathcal{M}_{\Lambda_{\mathbb{N}^+}^k}(u_X) \subseteq \mathcal{BE}_{\geq k}(u_X) \). On the other hand, \( \mathcal{BE}_{\geq k}(u_X) \subseteq \mathcal{M}_{\Lambda_{\mathbb{N}^+}^k}(u_X) \) by Theorem 3 and thus \( \mathcal{M}_{\Lambda_{\mathbb{N}^+}^k}(u_X) = \mathcal{BE}_{\geq k}(u_X) \).

Let us now suppose \( \mu \neq u_X \). Then, \( \exists A \neq X \) such that \( m(A) \neq 0 \). By Lemma 2, let us take \( A \) such that \( m(A) > 0, m(B) = 0, \forall B \subseteq A \). As \( A \neq X \), we can find \( j \notin A \). Consider \( i \in A \). We define \( \lambda \) by

\[
\lambda(A, \{i, j\}) = 1, \lambda(A, B) = 0 \text{ otherwise.}
\]

Now, if \( m(C) \neq 0 \), then \( C \not\subseteq A \) by hypothesis and thus, there exists \( l \in C \) such that \( l \notin A \). We define

\[
\lambda(C, l) = 1, \lambda(C, B) = 0 \text{ otherwise.}
\]

For \( C \subseteq A \), we complete \( \lambda \) in the conditions of Theorem 3. Then, \( m_\lambda(A) = 0 \) and it is \( m_\lambda(B) = 0, \forall B \subseteq A \), too. Therefore, \( \lambda_\mu(A) = 0 < \mu_\lambda(A) = m(A) \).

Consequently, we have \( \mathcal{BE}_{\geq k}(\mu) \subseteq \mathcal{M}_{\Lambda_{\mathbb{N}^+}^k}(\mu) \), whenever \( \mu \neq u_X \) and therefore, for \( k \)-additive dominating belief functions we have the same problems as for general \( k \)-additive measures.

At this point, it can be argued that in [1], good properties for Theorem 1 could be derived only when \( \mu \) is a belief function. Let us then study this case. If \( \mu \) is a belief function, the following can be proved:

**Proposition 11.** Consider a capacity \( \mu \). Then, \( \mathcal{M}_{\Lambda_{\mathbb{N}^+}^k}(\mu) \subseteq \mathcal{BE}_{k}(X) \) if and only if \( \mu \) is itself a belief function.

**Proof:** If \( \mu \) is a belief function, then by Lemma 3, \( \mathcal{M}_{\Lambda_{\mathbb{N}^+}^k}(\mu) \subseteq \mathcal{BE}_{k}(X) \).

Let us suppose now that \( \mu \) is not a belief function. Then, there exists \( A \) such that \( m_\mu(A) < 0 \). Consider \( i \in A \). We are going to prove that we can find a special choice of \( \lambda \in \Lambda_{\mathbb{N}^+}^k \) such that the condition \( m_\lambda(i) \geq 0 \) of Proposition 10 does not hold.

Let us take \( j \neq i \) and consider \( \lambda \) defined by:

\[
\lambda(A, i) = 1, \lambda(A, B) = 0 \text{ for } B \neq \{i\},
\]

\[
\lambda(C, i) = 0, \lambda(C, \{i, j\}) = 1, \forall C \supseteq i, C \neq A.
\]

Finally, for \( C \) such that \( i \notin C \), choose \( l \in C \) and define

\[
\lambda(C, l) = 1, \lambda(C, B) = 0 \text{ otherwise.}
\]

Then, \( \lambda \in \Lambda_{\mathbb{N}^+}^k \) and, for this choice of \( \lambda \), we have \( \mu_\lambda(i) = m_\lambda(i) = \sum_{B \ni i} \lambda(B, i)m(B) = m(A) < 0 \).

However, we cannot ensure the dominance as we have pointed in Proposition 10. On the other hand, we have

**Proposition 12.** [11] If \( \mu \) is a belief function and \( \lambda \in \Lambda_{\mathbb{N}^+}^k \), then the corresponding \( \mu_\lambda \) is a dominating belief function of \( \mu \).

As remarked in [11] and in Example 4, the set \( \mathcal{M}_{\Lambda_{\mathbb{N}^+}^k}(\mu) \) does not cope with all dominating belief functions and consequently, Proposition 8 cannot be extended. Then, if \( \mu \) is a belief function,

\[
\mathcal{M}_{\Lambda_{\mathbb{N}^+}^k}(\mu) \subseteq \mathcal{BE}_{k}(\mu) \subseteq \mathcal{M}_{\Lambda_{\mathbb{N}^+}^k}(\mu),
\]

with equality only if \( \mu = u_X \) by Proposition 10. Next result shows that all \( k \)-additive belief functions dominating another belief function \( \mu \) can be generated from \( \mathcal{M}_{\Lambda_{\mathbb{N}^+}^k}(\mu) \), the natural extension of \( \mathcal{M}_{\Lambda_{\mathbb{N}^+}^k}(\mu) \). This is the main result of the paper.
Theorem 4. Let \( \mu, m, \mu^*, m^* : \mathcal{P}(X) \to \mathbb{R} \), where \( \mu \) is a belief function, \( \mu^* \in \mathcal{BE}_{\Xi}(\mu) \) and \( m, m^* \) their corresponding Möbius inverses. Then, there exists \( \mu' \in \mathcal{BE}_{\Xi}(\mu) \) such that its Möbius transform \( m' \) can be written as

\[
m'(B) = \sum_{A \mid B \subseteq A} \lambda'(A, B)m(A), \forall B \in \mathcal{P}^k(X),
\]

where \( \lambda' : \mathcal{P}(X) \times \mathcal{P}^k(X) \to [0, 1] \) is such that

\[
\sum_{B \mid B \subseteq A} \lambda'(A, B) = 1, \forall A \in \mathcal{P}(X).
\]  

(5)

\[
\lambda'(A, B) = 0 \text{ if } B \nsubseteq A,
\]

(6)

and \( m^* \) can be derived from \( m' \) through

\[
m^*(C) = \sum_{B \mid B \subseteq C} \lambda^*(B, C)m'(B), \forall C \in \mathcal{P}^k(X),
\]

(7)

where \( \lambda^* : \mathcal{P}^k(X) \times \mathcal{P}^k(X) \to [0, 1] \) is such that

\[
\sum_{C \mid B \subseteq C} \lambda^*(B, C) = 1, \forall B \in \mathcal{P}^k(X).
\]

(8)

\[
\lambda^*(B, C) = 0 \text{ if } B \nsubseteq C.
\]

(9)

Proof: This result is explained in Figure 2 for \( |X| = 3 \) and \( k = 2 \).

![Figure 2: Example of flow network for \(|X| = 3 \) and \( k = 2 \).](image)

We already know from Theorem 3 that

\[
m^*(C) = \sum_{A \mid C \cap A \neq \emptyset} \lambda(A, C)m(A), \lambda(A, C) \geq 0, \sum_{A \mid A \cap C \neq \emptyset} \lambda(A, C) = 1, \forall A \in \mathcal{P}(X), \forall C \in \mathcal{P}^k(X).
\]

Let us define a weight function \( \lambda' \) by

\[
\lambda'(A, B) := \sum_{C \mid C \cap A = B} \lambda(A, C), \forall A \in \mathcal{P}(X), \forall B \in \mathcal{P}^k(X), B \subseteq A, \lambda'(A, B) = 0 \text{ otherwise.}
\]

(10)
Then, we define \( m' \) by

\[
m'(B) := \sum_{A \subseteq B} \lambda'(A, B)m(A), \ \forall B \in \mathcal{P}^k(X). \tag{11}
\]

Let us check that \( \lambda' \) satisfies the conditions imposed in the theorem.

- \( \lambda'(A, B) = \sum_{C \subseteq A, C \cap A = B} \lambda(A, C) \geq 0, \ \forall A \in \mathcal{P}(X), \ \forall B \in \mathcal{P}^k(X), \) as \( \lambda(A, C) \geq 0, \ \forall A \in \mathcal{P}(X), \ \forall C \in \mathcal{P}^k(X) \) by Theorem 3.

- \( \sum_{B \subseteq A} \lambda'(A, B) = \sum_{B \subseteq A} \lambda(A, C) = \lambda(A, C) = 1, \ \forall A \in \mathcal{P}(X) \) by Theorem 3, so that Equation (5) holds.

- Equation (6) is satisfied by the construction of \( \lambda' \).

Let us now prove that \( \mu' \) is a belief function: First,

\[
m'(B) = \sum_{A \subseteq B} \lambda'(A, B)m(A) \geq 0, \forall B \in \mathcal{P}^k(X)
\]

because \( \mu \) is a belief function (and then \( m(A) \geq 0, \ \forall A \subseteq X \) by Proposition 2), and \( \lambda'(A, B) \geq 0, \ \forall A \in \mathcal{P}(X), \ \forall B \in \mathcal{P}^k(X) \).

Then, it suffices to prove that \( \sum_{B \subseteq X} m'(B) = 1 \), but this holds since

\[
\sum_{B \subseteq X} m'(B) = \sum_{B \subseteq X} \lambda'(A, B)m(A) = \sum_{A \subseteq X} \sum_{B \subseteq A} \lambda'(A, B)m(A) = \sum_{A \subseteq X} m(A) = 1.
\]

- Besides, \( \mu' \) dominates \( \mu \):

\[
\mu'(B) = \sum_{D \subseteq B} m'(D) = \sum_{D \subseteq B} \lambda'(A, D)m(A).
\]

If \( A \subseteq B \), we have that whenever \( D \subseteq A \), then \( D \subseteq B \), and hence, \( m(A) \) is multiplied by

\[
\sum_{D : D \subseteq A} \lambda'(A, D) = 1.
\]

Consequently,

\[
\sum_{D \subseteq B} m'(D) = \sum_{A \subseteq B} m(A) + \sum_{A \subseteq B} \lambda'(A, D)m(A).
\]

Finally, as \( m(A) \geq 0, \ \forall A \in \mathcal{P}(X), \ \lambda'(A, B) \geq 0, \ \forall A \in \mathcal{P}(X), \ \forall B \in \mathcal{P}^k(X) \), we obtain that

\[
\mu'(B) \geq \sum_{A \subseteq B} m(A) = \mu(B).
\]

- \( \mu' \) generates \( \mu^* \) through Equation (7): For each \( B, C \) such that \( B \subseteq C \) let us define

\[
\lambda^*(B, C) := \sum_{A : A \cap C = B} \lambda(A, C) \frac{m(A)}{m'(B)}, \ \text{if } m'(B) \neq 0.
\]

What happens if \( m'(B) = 0 \)? In this case, for all \( A \) such that \( B \subseteq A \), we have \( \lambda'(A, B) = 0 \) or \( m(A) = 0 \) by Equation (11) and thus, \( \lambda(A, C) = 0 \), or \( m(A) = 0, \ \forall C \) such that \( A \cap C = B \) by Equation (10). We can then choose arbitrary values for \( \lambda^*(B, C) \) under the conditions of Equations (8) and (9).
Then,
\[
\sum_{B \subseteq C} \lambda^*(B, C) m'(B) = \sum_{B \subseteq C, A \cap C = B} \lambda(A, C) \frac{m(A)}{m'(B)} m'(B) = \sum_{B \subseteq C, A \cap C = B} \lambda(A, C) m(A) = \sum_{A \mid A \cap C \neq \emptyset} \lambda(A, C) m(A) = m^*(C),
\]
by Theorem 3. Thus, with our definitions of \(\mu'\) and \(\lambda^*\), we have obtained \(\mu^*\).

It is clear that \(\lambda^*(B, C) \geq 0, \forall B, C \in \mathcal{P}^k(X)\) as \(\lambda(A, C) \geq 0, m(A) \geq 0, m'(B) \geq 0, \forall A \in \mathcal{P}(X), \forall B \in \mathcal{P}^k(X)\).

To finish the proof, it suffices to show that for any \(B \in \mathcal{P}^k(X)\), we have \(\sum_{C \mid B \subseteq C} \lambda^*(B, C) = 1\). We have
\[
\sum_{C \mid B \subseteq C} \lambda^*(B, C) = \sum_{C \mid B \subseteq C} \sum_{A \cap C = B} \lambda(A, C) \frac{m(A)}{m'(B)} = \frac{1}{m'(B)} \sum_{A \mid C \cap A = B} \lambda(A, C) m(A) = \frac{1}{m'(B)} \sum_{A \subseteq X \mid B \subseteq A} \lambda(A, B) m(A) = \frac{1}{m'(B)} m'(B) = 1.
\]

Then, the result holds.
\[
\square
\]

As a final remark, note that we can find some examples where we can ensure \(\lambda \geq 0\), even if \(\mu\) is not a belief function:

**Example 5.** Consider \(|X| = 4\) and the capacity defined by \(\mu(A) = 1\), if \(A \neq \emptyset\) (this is \(u_0\), which is not a belief function by Proposition 2). Then, the only \(n\)-additive capacity that dominates \(\mu\) is \(\mu\) itself. Then, function \(\lambda\) that determines the only capacity dominating \(\mu\) (i.e. \(\mu\)) is given by \(\lambda(A, A) = 1, \lambda(A, B) = 0\) if \(A \neq B\), and thus \(\lambda\) is a non-negative function, even if \(\mu\) is not a belief function.

5 Application to decision making

We illustrate in this section several possible applications of the above results, i.e., the construction of the set of \(k\)-additive belief functions dominating a given capacity.

A whole set of applications comes from the following situation: suppose that we want to build a decision model from some elicitation procedure with a human decision maker. Most of the time, the decision maker is not able to provide enough information or enough precise answers in order to have a unique possible decision model compatible with the information obtained from the decision maker. In our case, we make the assumption that our decision model is based on capacities, which is known in many fields of decision making to provide versatile models. This means that the parameters of our model are precisely the coefficients of a capacity \(\mu\), and that we cannot expect better information than \(\mu(A)\) belongs to some interval \([\mu_+(A), \mu^+(A)]\), for all subset \(A\) of the referential set \(X\).

Let us first consider the lower bound \(\mu_+\), since in many cases the upper bound may be simply the uninformative value 1. Depending on the specific domain of application and type of decision problem, one may want to restrict to some particular family of capacities, say \(\mathcal{F}\), so that the problem of identification of the decision model writes: find all capacities in \(\mathcal{F}\) dominating \(\mu_+\). The choice of a particular family of capacities is principally based on two reasons:

- we are looking for capacities satisfying a given set of properties, most often properties linked to decision behaviour (attitude towards risk, etc.). This is related to the axiomatic approach to decision making.
• to limit the number of possible solutions and to limit complexity of the model, we impose some restrictions, typically, $k$-additivity.

The case of $k$-additive belief functions appear as a natural choice for $F$ at least in the following cases:

• in social welfare theory, the case of $k$-additive belief functions has been characterized by the authors [22], following previous works of Ben Porath and Gilboa [23], and Gajdos [7], with axioms which are natural in the context of social welfare since they are closely related to the well-known Gini index. In particular, it is proved that $k$-additivity can be seen as an aversion to social inequality.

• in decision under uncertainty or risk, the meaning of a belief function, compared to a classical probability measure, is well known. Roughly speaking, belief functions permit to distinguish between ignorance and equiprobability, a feature which is particularly useful for handling missing or ambiguous data. There also exist various axiomatic characterization results of belief functions, see e.g. Jaffray and Wakker [14], and Wakker [29].

For example, assume $|X| = 3$ and suppose that from the information given by the decision maker, we can only deduce that $\mu_\ast(3) \geq 0.4$ and $\mu_\ast(A) \in [0, 1]$ for $A \neq \{3\}$. Then, the lower measure is a belief function given by $m_\ast(3) = 0.4, m_\ast(X) = 0.6, m_\ast(A) = 0$, otherwise.

Then, the set of all 2-additive belief measures compatible with the information is given by the set of dominant belief functions in $M_{\Lambda_2}^\ast(\mu)$ by Theorem 3, and these measures can be generated from $M_{\Lambda_2}^\ast,+,+ (\mu)$ by Theorem 4.

In order to be complete, the upper bound should be considered as well, which motivates a companion study for the set of $k$-additive belief functions dominated by some capacity.

A second possible application arises in cooperative game theory, where $X$ represents the set of players. The notion of dominating probability measure (in fact additive non normalized capacities since in this context $\mu(X)$ represents the total worth of the game, which in general is not equal to 1) is known in this domain under the name of “core of a game”. The existence of the core is closely related to the formation of coalitions, since any additive (non normalized) capacity in the core represents a distribution of some non negative amount to each player. Defining the core as the set of dominating $k$-additive belief functions generalizes the classical definition in the sense that the distribution will be defined over individual players and all coalitions of at most $k$ players.

The fact that we impose belief functions guarantees that the amount will be non negative, since it corresponds to the Möbius transform.

In all cases, the notion of $k$-additive belief function appears as the most natural generalization of probability measures, the value of $k$ measuring the complexity of the generalization. In this finite setting, a probability measure is defined by its probability distribution over all singletons, which is a non-negative amount. A 2-additive belief function induces a non negative distribution over singletons and pairs, and so on.

6 Conclusions and open problems

In this paper we have studied the set of $k$-additive belief functions dominating a given capacity. This set is a middle term between the set of dominating probabilities and the set of dominating $k$-additive measures.

We have seen that the results for the general $k$-additive case do not allow a translation of the results derived for probabilities in [1]. This is due to the fact that the conditions are too general, thus leading a too wide set of functions.

We have obtained a result, similar to those for dominating probabilities and $k$-additive measures, that gives us a family of functions in which the set of dominating $k$-additive belief functions is included. Our results for belief functions permit a reduction in the set of functions. However, the family obtained in this case is not the natural extension of the corresponding family obtained for probabilities and provides again a too wide set of functions.

Finally, we have proved that it suffices to consider this natural extension in order to derive any other dominating $k$-additive belief function whenever the initial measure is a belief function, too.

We have proposed several possible applications in decision making, especially in social welfare and decision under risk or uncertainty.
As immediate research, we have the problem of finding the profile of $BE\mathcal{L}^{\infty}_{\mu}(\mu)$ when $\mu$ is a 2-monotone measure or a belief function. In this sense, we have already found some results in [21], where Theorem 4 plays an important role.

Another interesting problem is to look for extensions of Theorem 4. Finally, two other questions arise from a theoretical point of view:

1. We have proved that for belief functions, $\lambda \geq 0$. However, this is not a necessary condition (Example 5). Then, we can ask ourselves for necessary and sufficient conditions on $\mu$ for $\lambda \geq 0$.

2. The same problem can be considered for the $A \subseteq B$ condition.

7 Acknowledgements

This research has been supported in part by FEDER-MCYT grant number BFM2001-3515.

References


